

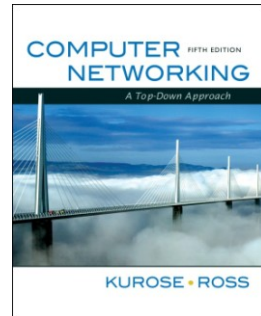
RSC

Part I: Introduction

Redes y Servicios de Comunicaciones Universidad Carlos III de Madrid

These slides are, mainly, part of the companion slides to the book "Computer Networking: A Top Down Approach" generously made available by their authors (see copyright below). The slides have been adapted, where required, to the teaching needs of the subject above.

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*Computer Networking:
A Top Down Approach
5th edition.*

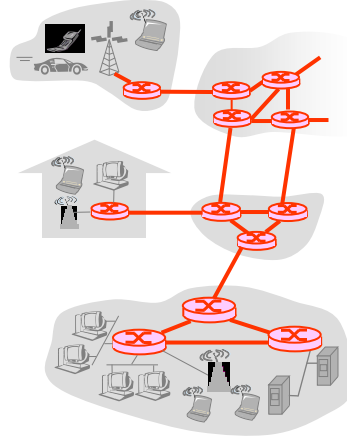
Jim Kurose, Keith Ross
Addison-Wesley, April
2009.

RSC Part I: Introduction

- ❑ Circuit switching vs packet switching
- ❑ Protocols and protocols stacks
- ❑ What is the Internet
- ❑ Network structure
- ❑ ISPs and Internet Backbones

The Network Core

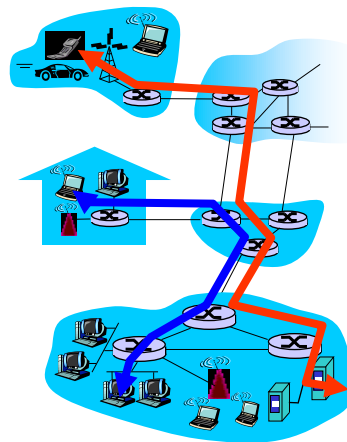
- ❑ mesh of interconnected routers
- ❑ the fundamental question: how is data transferred through net?
 - ❖ **circuit switching**: dedicated circuit per call: telephone net
 - ❖ **packet-switching**: data sent through net in discrete "chunks"



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Network Core: Circuit Switching

- End-end resources reserved for "call"**
- ❑ link bandwidth, switch capacity
 - ❑ dedicated resources: no sharing
 - ❑ circuit-like (guaranteed) performance
 - ❑ call setup required



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Network Core: Circuit Switching

network resources
(e.g., bandwidth)

divided into "pieces"

- pieces allocated to calls
- resource piece *idle* if not used by owning call (*no sharing*)

□ dividing link bandwidth into "pieces"

- ❖ frequency division
- ❖ time division

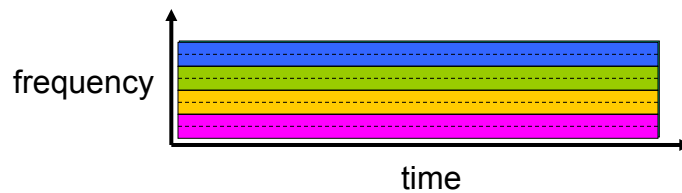
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Circuit Switching: FDM and TDM

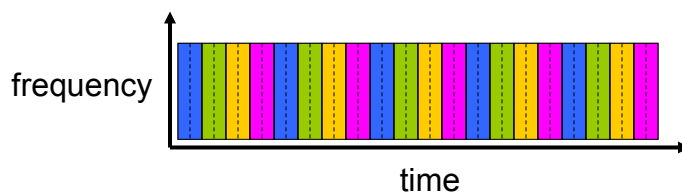
FDM

Example:

4 users



TDM



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Network Core: Packet Switching

each end-end data stream
divided into *packets*

- ❑ user A, B packets *share* network resources
- ❑ each packet uses full link bandwidth
- ❑ resources used *as needed*

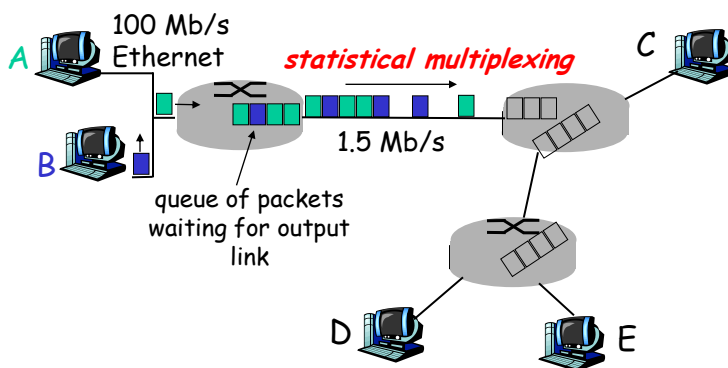
Bandwidth division into "pieces"
Dedicated allocation
Resource reservation

resource contention:

- ❑ aggregate resource demand can exceed amount available
- ❑ congestion: packets queue, wait for link use
- ❑ store and forward: packets move one hop at a time
 - ❖ Node receives complete packet before forwarding

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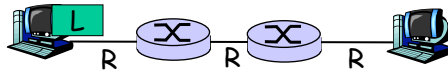
Packet Switching: Statistical Multiplexing



Sequence of A & B packets does not have fixed pattern,
bandwidth shared on demand → **statistical multiplexing**.
TDM: each host gets same slot in revolving TDM frame.

Introduction 1-8

Packet-switching: store-and-forward



- ❑ takes L/R seconds to transmit (push out) packet of L bits on to link at R bps
- ❑ *store and forward*: entire packet must arrive at router before it can be transmitted on next link
- ❑ delay = $3L/R$ (assuming zero propagation delay)

Example:

- ❑ $L = 7.5$ Mbits
- ❑ $R = 1.5$ Mbps
- ❑ transmission delay = 15 sec

} more complex than this ...

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Packet switching versus circuit switching

Packet switching allows more users to use network!

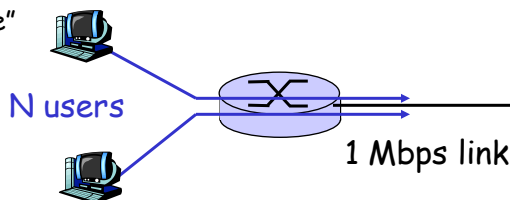
- ❑ 1 Mb/s link
- ❑ each user:
 - ❖ 100 kb/s when "active"
 - ❖ active 10% of time

circuit-switching:

- ❖ 10 users

packet switching:

- ❖ with 35 users, probability > 10 active at same time is less than .0004



Q: how did we get value 0.0004?

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Packet switching versus circuit switching

Is packet switching a "slam dunk winner?"

- ❑ great for bursty data
 - ❖ resource sharing
 - ❖ simpler, no call setup
- ❑ **excessive congestion:** packet delay and loss
 - ❖ protocols needed for reliable data transfer, congestion control
- ❑ **Q: How to provide circuit-like behavior?**
 - ❖ bandwidth guarantees needed for audio/video apps
 - ❖ still an unsolved problem

Introduction 1-11

What's a protocol?

human protocols:

- ❑ "what's the time?"
- ❑ "I have a question"
- ❑ introductions

... specific msgs sent

... specific actions taken
when msgs received,
or other events

network protocols:

- ❑ machines rather than humans
- ❑ all communication activity in Internet governed by protocols

*protocols define format,
order of msgs sent and
received among network
entities, and actions
taken on msg
transmission, receipt*

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Protocol "Layers"

Networks are complex!

- ❑ many "pieces":
 - ❖ hosts
 - ❖ routers
 - ❖ links of various media
 - ❖ applications
 - ❖ protocols
 - ❖ hardware, software

Question:

Is there any hope of
organizing structure of
network?

Or at least our discussion
of networks?

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Why layering?

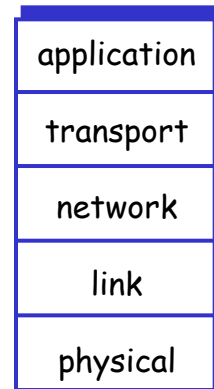
Dealing with complex systems:

- ❑ explicit structure allows identification, relationship of complex system's pieces
 - ❖ layered **reference model** for discussion
- ❑ modularization eases maintenance, updating of system
 - ❖ change of implementation of layer's service transparent to rest of system
 - ❖ e.g., change in gate procedure doesn't affect rest of system
- ❑ layering considered harmful?

Introduction 1-14

Internet protocol stack

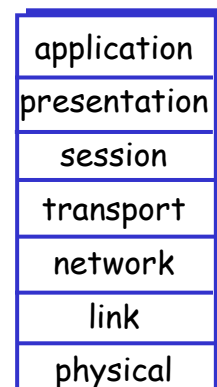
- ❑ **application:** supporting network applications
 - ❖ FTP, SMTP, HTTP
- ❑ **transport:** process-process data transfer
 - ❖ TCP, UDP
- ❑ **network:** routing of datagrams from source to destination
 - ❖ IP, routing protocols
- ❑ **link:** data transfer between neighboring network elements
 - ❖ PPP, Ethernet
- ❑ **physical:** bits "on the wire"



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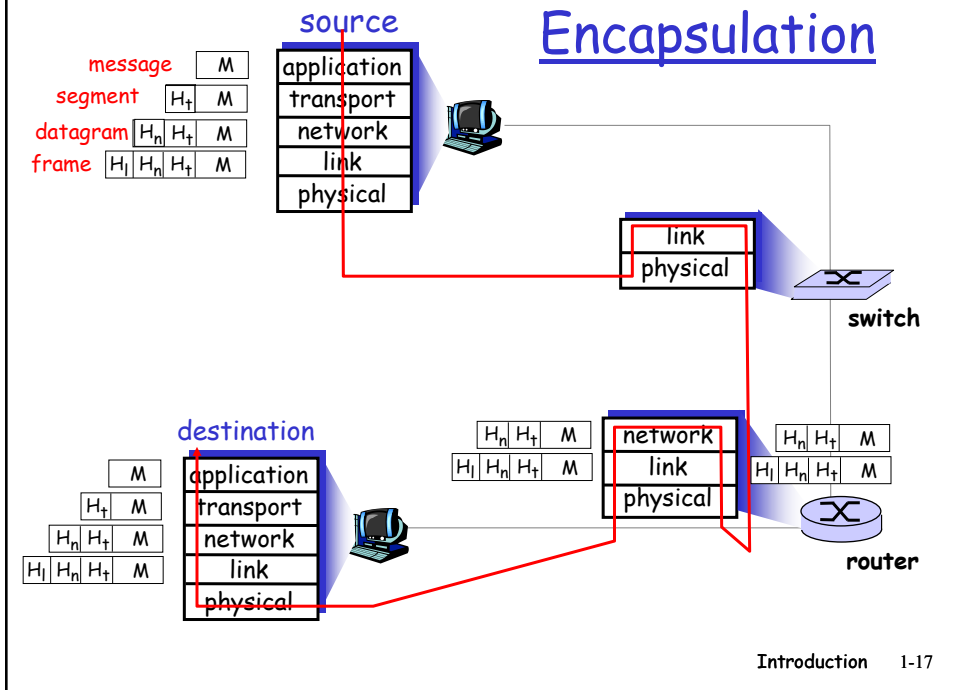
ISO/OSI reference model

- ❑ **presentation:** allow applications to interpret meaning of data, e.g., encryption, compression, machine-specific conventions
- ❑ **session:** synchronization, checkpointing, recovery of data exchange
- ❑ Internet stack "missing" these layers!
 - ❖ these services, *if needed*, must be implemented in application
 - ❖ needed?



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Encapsulation



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